

Opportunistic Spectrum Aggregation for Cognitive Communications Under Collision Constraints

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Abstract—We consider a collision-sensitive secondary system that intends to opportunistically aggregate and utilize spectrum of a primary system to achieve higher data rates. In such opportunistic spectrum access, secondary transmission can collide with primary transmission. When the secondary system aggregates more channels for data transmission, more frequent collisions may occur, limiting the performance obtained by the opportunistic spectrum aggregation. In this context, dynamic spectrum aggregation problem is formulated to maximize the ergodic channel capacity under the constraint of collision tolerable level. To solve the problem, we develop the optimal spectrum aggregation approach, deriving closed-form expressions for the collision probability in terms of primary user traffic load, secondary user transmission interval, and the random number of sub-channels aggregated. Our results show that aggregating only a subset of sub-channels will be a better choice, depending on the ratio of collision sensitivity requirement to the primary user traffic.

Keywords—Spectrum aggregation, Collision probability, Opportunistic Spectrum Access

I. INTRODUCTION

By enabling the secondary users (SUs) to exploit the unoccupied spectrum of the primary users (PUs), the opportunistic spectrum access is regarded as a promising solution for resolving spectrum underutilization [1].

In the opportunistic spectrum access, since PUs have the priority over SUs in accessing the channels, PUs can access the channel a SU is using at any time. If such collision happens between PU and SU transmission, the SU has to vacate the channel immediately and move to a new one. While the SU communication has to be interrupted, packets must wait in the transmission buffer. The communication can be resumed when a connection is successfully established on a new channel. Such spectrum handoff naturally causes additional latency that affects SU performance in addition to causing short-term interference to PUs [2].

When a SU requires high data rates, the SU can use multiple channels simultaneously through spectrum aggregation [3]. However, the use of an aggregate channel comprising of multiple channels could incur frequent collisions. While improvement of the SU's performance could be expected by spectrum aggregation, frequent collisions would result in performance degradation of the SU. It is possible that the SU receives little benefit from spectrum aggregation. Thus, when the SU aggregates multiple channels, collision issues should be considered more carefully.

However, in literature, spectrum aggregation strategies of the secondary system proposed to date, only aim to improve the spectrum utilization without consideration of collision. In [4][5], the research focuses on how many channels should be aggregated, in order to accommodate more SUs and to improve data rates. In [6][7], while considering the hardware constraint for spectrum aggregation, aggregation algorithms with only the aim to achieve higher network throughputs have been proposed. Even though the collision probability of the SU [8] and spectrum allocation strategies for the secondary system considering the SU's collision probability have been investigated in [9][10][11], they only consider the single channel use by the SU. To the best of our knowledge, however, spectrum aggregation by secondary users considering the collision issue has not been covered in the literature.

In this paper, we consider a collision-sensitive secondary system opportunistically accessing multiple primary channels and aim at developing a new method for dynamic spectrum aggregation so that the ergodic capacity of an aggregate channel is maximized under the constraint that the collision probability is guaranteed to be smaller than some threshold value. For the aggregate channel use, the closed-form expression of the collision probability is derived. By analysing the impact of spectrum aggregation on the ergodic capacity and the collision probability, we address the challenge to maximize the ergodic capacity of the aggregate channel under the collision probability constraint. Moreover, a new optimization problem is formulated and by solving it, we intend to answer to a fundamental question, how many sub-channels should be activated for the spectrum aggregation.

The rest of this paper is organized as follows. Section II presents the system and channel models for primary and secondary system, describing the traffic pattern and the dynamic spectrum aggregation scheme. The spectrum aggregation problem to maximize the ergodic channel capacity considering the collision probability is formulated in Section III. The ergodic capacity and the collision probability are analysed in Section IV. Then, the optimal solution is addressed in Section V. After showing the numerical and simulation results in Section VI, our conclusions are summarized in Section VII.

II. SYSTEM AND CHANNEL MODEL

A. System model

We consider a secondary system coexisting with a primary system. The primary system communicates with the PUs

through multiple licensed channels. The PUs, as the licensed users, have priority over the SUs in accessing the licensed (primary) channels. Thus, the secondary system opportunistically detects and aggregates idle primary channels, and allocates the aggregate channel (consisting of multiple primary channels) for the data transmission to the SU.

We assume that in the primary system, the packet arrival rate at each channel follows a Poisson process with a rate λ_p . The service time of the packets, $1/\mu_p$, has the general distribution [11]. The packet arrival rate and service time at each channel are assumed to be independent and identically distributed (i.i.d.).

Suppose that in the secondary system, the base station (BS) is able to identify the availability of primary channels by spectrum sensing without errors. Particularly, every sensing period, the BS will identify a subset of primary channels unoccupied by PUs. For given subset of N unoccupied channels, the BS sends the data to a SU during the data transmission interval. The data transmission interval of the secondary system is denoted by T_d , which is in general much longer than the sensing interval. T_d is the maximum time interval over which a SU remains unaware of any changes in channel occupancy.

Consider a realistic case of the opportunistic spectrum access where PUs can return to the channels that are still used by the SU. In such case, collision can happen between primary and secondary transmissions. Such collision is inherent in realistic systems where *no synchronization* between the primary and secondary transmissions is assumed. To measure such collision, we define the collision probability, P_c , as the probability that the secondary transmission collides with the primary transmission and it can be given by [8]

$$P_c = \lim_{T \rightarrow \infty} \frac{\text{the number of collided SU transmissions in } [0, T]}{\text{the number of SU transmissions in } [0, T]}. \quad (1)$$

In the presence of certain collision, the secondary system can perform effective communication while P_c remains below the pre-defined threshold of collision (denoted by ξ_{th}). The collision threshold indicates the maximum collision probability tolerated by the secondary system. Notice that the value of ξ_{th} can influence the performance of the secondary system as well as that of the primary system. That is, until PUs are discovered from the time they collide with the SU, PUs can suffer from short-term interference. Thus, the quality of service of PUs could be also considered in the setting of ξ_{th} .

B. Channel model

For given N unoccupied primary channels, we propose that the BS of the secondary system aggregates only a subset of n primary channels, properly selecting the value of $n (\leq N)$. Particularly, the BS randomly selects n among N channels every transmission interval T_d . That is, each one of the N channels is *equally likely* to be selected and allocated at every transmission interval. Let the maximum transmission power be denoted by P_{max} , and P_{max} is *equally distributed* among the n randomly selected channels. Since the n primary channels now comprise an aggregate channel, hereafter, the term ‘primary channel’ can be interchangeable with ‘sub-channel (of an aggregate channel)’.

Suppose that available sub-channels between the BS and the SU are independent and Rayleigh flat fading. The channel coefficient of sub-channel i (denoted by h_i , for all i) is complex Gaussian random variable with zero mean and variance $\sigma_{h_i}^2$, i.e., $h_i \sim \mathcal{CN}(0, \sigma_{h_i}^2)$. Such channel state information (CSI) is assumed to be perfectly known at the receiver. Then, the received signal to noise ratio (SNR) at sub-channel i for $i \in \{1, 2, \dots, N\}$ can be represented as

$$\rho_i = \frac{|h_i|^2 P_i}{\sigma^2} = \frac{g_i P_i}{\sigma^2}, \quad (2)$$

where $g_i = |h_i|^2, \forall i$, is a Chi-square distributed random variable with $2k$ degrees of freedom where k denotes the number of receiver antennas along with multi-antenna techniques (e.g., maximal ratio combining) [12]. P_i is the transmit power, i.e., $P_i = P_{max}/n$ for equal power allocation. Notice in (2) that σ^2 is the variance of the complex-valued zero-mean additive white Gaussian noise (AWGN).

We can express the capacity of sub-channel i of bandwidth B as $C_i = B \log_2(1 + \rho_i)$ [13]. Using (2), C_i can be given by

$$C_i(n, \frac{g_i}{\sigma^2}) = B \cdot \log_2 \left(1 + \frac{g_i \cdot P_{max}}{n \cdot \sigma^2} \right). \quad (3)$$

From (3), the ergodic capacity for the selected sub-channel i can be obtained by taking the expectation of (3) with respect to g_i/σ^2 which can be expressed as

$$E[C_i] = \int_0^\infty B \cdot \log_2 \left(1 + P_{max} \frac{x}{n} \right) p_x(x) dx, \quad (4)$$

where $E[\cdot]$ stands for the expectation operator and $p_x(x)$ is the pdf of $x = g_i/\sigma^2$, for all i , which can be given by

$$p_x(x) = \frac{1}{2^{k/2} \Gamma(\frac{k}{2})} x^{\frac{k}{2}-1} e^{-\frac{x}{2}},$$

where $\Gamma(\cdot)$ denotes the gamma function [13].

The capacity of the aggregate channel is upper bounded by the sum of the capacity of the n active sub-channels, as the BS aggregates n sub-channels to transmit the data. For the sake of the simplicity, let B be normalized (i.e., $B = 1$). Due to the proposed aggregation rule above, each of the N sub-channels has the same probability to be randomly selected, according to the uniform distribution. The ergodic capacity, C_{total} , of the aggregate channel is calculated, weighting the ergodic capacity in (4) by its own aggregation probability and adding them all. Thus, C_{total} can be expressed as

$$C_{total} = \frac{n}{N} \sum_{i=1}^N E \left[\log_2 \left(1 + \frac{g_i \cdot P_{max}}{n \cdot \sigma^2} \right) \right]. \quad (5)$$

III. PROBLEM FORMULATION

We study the problem of the spectrum aggregation, properly selecting only a subset of n among N total idle sub-channels for the secondary transmission. This is to maximize the achievable data rate C_{total} while satisfying a collision probability requirement. Using (5), this problem can be posed as

$$\begin{aligned} \max C_{total} &= \max \frac{n}{N} \sum_{i=1}^N E \left[\log_2 \left(1 + \frac{g_i \cdot P_{max}}{n \cdot \sigma^2} \right) \right], \\ &\text{subject to } P_{c,n} \leq \xi_{th}. \end{aligned} \quad (6)$$

where $P_{c,n}$ denotes the probability that secondary transmission exploiting the n sub-channels collides with the primary transmission.

Notice that collision by the secondary transmission on the n sub-channels includes all the events of collision across all possible subsets of the n sub-channels. Thus, for a given n , $P_{c,n}$ can be derived, utilising the collision probability of an individual sub-channel as follows:

$$\begin{aligned} P_{c,n} &= \text{P}\{\text{Collision in any of } n \text{ channels}\} \\ &= 1 - \text{P}\{\text{No collision in all } n \text{ channels}\}. \end{aligned} \quad (8)$$

To solve the problems (6)-(8), we raise and address a fundamental question, ‘‘how many sub-channels must be aggregated for higher data rate under the collision requirement?’’ To that end, we first start by analysing the impact of the number of aggregate sub-channels on the performance, followed by the optimal solution using analytical results.

IV. ANALYSIS ON ERGODIC CAPACITY AND COLLISION PROBABILITY

A. Impact of the channel aggregation on C_{total}

Since the log function is a concave function, by using Jensen’s inequality [13], the ergodic capacity of the aggregate channel composed of the n sub-channels for a given N can be obtained by

$$\begin{aligned} C_{total} &= \frac{n}{N} \sum_{i=1}^N E \left[\log_2 \left(1 + \frac{g_i \cdot P_{max}}{n \cdot \sigma^2} \right) \right] \\ &\leq \frac{n}{N} \sum_{i=1}^N \log_2 \left(1 + \frac{\gamma_i}{n} \right), \end{aligned} \quad (9)$$

where γ_i denotes the average received SNR level with the maximum transmit power on sub-channel i , i.e., $\gamma_i = E[g_i] \cdot P_{max} / \sigma^2$.

Let us figure out how the ergodic capacity behaves with n , the number of sub-channels used for aggregation. For this, the first and the second derivatives of the ergodic capacity with respect to n are computed as follows

$$\frac{\partial C_{total}}{\partial n} = \frac{1}{N(\log 2)} \sum_{i=1}^N \left[\log \left(1 + \frac{\gamma_i}{n} \right) - \frac{\gamma_i}{n + \gamma_i} \right]. \quad (10)$$

$$\frac{\partial^2 C_{total}}{\partial^2 n} = \frac{1}{N(\log 2)} \sum_{i=1}^N \frac{-\gamma_i^2}{n(\gamma_i + n)^2}. \quad (11)$$

Consider the homogeneous channel setup where $\gamma_i = \gamma, \forall i$. Noticing the fact that $\partial^2 C_{total} / \partial^2 n$ in (11) is not positive for all possible n , C_{total} is a concave function of n . The first derivative function is always positive (i.e., $\partial C_{total} / \partial n > 0$). This reveals that the ergodic capacity monotonically increases with n . Therefore, for a given finite N , *more n , higher C_{total} is obtained.*

B. Impact of the channel aggregation on collision probability

Let us consider the homogeneous case when the PU traffic intensity on sub-channel $i, \forall i$ are i.i.d. The collision probability of each sub-channel is equally likely being P_c in (1). Using this, (8) can be further simplified to

$$P_{c,n} = 1 - \prod_{i=1}^n (1 - P_c) = 1 - (1 - P_c)^n. \quad (12)$$

It can be shown in (12) that the collision probability, $P_{c,n}$ can be derived as an increasing function of the number of sub-channels, n , once the collision probability by a single sub-channel use, P_c is given.

To further analysis on $P_{c,n}$, we now calculate P_c . The collision happens only if PUs reappear on the sub-channels on which the SU is still transmitting. As shown in [8], the collision probability of a single sub-channel due to asynchronous transmission between primary and secondary transmission, P_c , is equivalent to the probability that at least one PU’s packet arrives during a SU’s transmission period T_d . As PU’s packets arrive according to a Poisson process, P_c can be given by

$$P_c = 1 - \exp(-\lambda_P \cdot T_d). \quad (13)$$

From (12) and (13), $P_{c,n}$ can be given, for given n aggregated sub-channels, by

$$P_{c,n} = 1 - (1 - P_c)^n = 1 - \exp(-\lambda_P \cdot T_d \cdot n). \quad (14)$$

In (14), $P_{c,n}$ indicates the probability that at least one primary packet arrives during a secondary transmission period in n sub-channels. It can be seen from (14) that $P_{c,n}$ increases with the product of λ_P and T_d , for given n aggregated sub-channels. For given λ_P and T_d , also, it can be shown from (14) that $P_{c,n}$ increases significantly with n .

V. OPTIMAL SOLUTION

We consider the collision-sensitive secondary system where it is desired to provide a realistic optimal solution that maximizes the ergodic capacity C_{total} , while simultaneously limiting the collision probability $P_{c,n}$ below the desired level. For this, we solve the maximization problem described in (6) and (7).

As per C_{total} and $P_{c,n}$, the analytical results in Section IV presented that the both are functions of the number n of active sub-channels used for aggregation. Particularly, notice the fact that the monotonicity of C_{total} with n presented in Section IV. Then, it can be stated that, for a given collision probability threshold ξ_{th} , the optimal value of $n(\leq N)$ maximizing C_{total} may exist and can be found as the *largest* among possible integer values satisfying the collision probability requirement ($P_{c,n} \leq \xi_{th}$). The requirement (7) can be rewritten, using (12), as

$$1 - (1 - P_c)^n \leq \xi_{th}. \quad (15)$$

Taking the logarithm of both sides in (15), the necessary condition on being possible candidate ns for the optimum can be obtained for a given ξ_{th} as

$$n \leq \left\lceil \frac{\log(1 - \xi_{th})}{\log(1 - P_c)} \right\rceil \quad (16)$$

where $\lfloor x \rfloor$ denotes the largest integer not larger than x , and P_c denotes the collision probability being (13). Here, note that n is the number of active sub-channels for aggregation, and thus the valid range of n in (16) must be given to $1 \leq n \leq N$.

Considering the requirement in (15), the optimum n^* must be the largest *integer* satisfying the inequality of (16). Thus, n^* subject to $P_{c,n} \leq \xi_{th}$ can be expressed using (16).

$$n^* = \begin{cases} \left\lfloor \frac{\log(1-\xi_{th})}{\log(1-P_c)} \right\rfloor, & \text{if } 1 \leq \left\lfloor \frac{\log(1-\xi_{th})}{\log(1-P_c)} \right\rfloor \leq N \\ N, & \text{elseif } \left\lfloor \frac{\log(1-\xi_{th})}{\log(1-P_c)} \right\rfloor > N \\ 0, & \text{otherwise} \end{cases}. \quad (17)$$

It can be shown from (17) that n^* relies on both P_c and ξ_{th} for a given N .

To further specify n^* , P_c in (13) are now taken into account, solving (17). To that end, we need to derive the expression for n^* with respect to the primary packet arrival rate λ_p and the SU transmission interval T_d . For simplicity in analysis, hereinafter, we consider only when all the candidates satisfying (16) are below N , i.e., $1 \leq \lfloor x \rfloor \leq N$, for a given N . This leads us to consider only the first inequality condition in (17). Since in the other conditions in (17) the corresponding optimal solutions can be provided straightforwardly, being beyond our interests.

Inserting (13) into (17), therefore, n^* can be found as follows

$$n^* = \left\lfloor \frac{\log(1-\xi_{th})}{-\lambda_p \cdot T_d} \right\rfloor. \quad (18)$$

In (18), it is observed that for a given collision requirement ξ_{th} , the optimal number of sub-channel for aggregation should be chosen inversely with a product of $\lambda_p \cdot T_d$. Particularly, it can be found from (18) that lower λ_p (or smaller T_d), more n^* can be. This reveals that *as the product $\lambda_p \cdot T_d$ decreases, the optimum number n^* of the aggregate sub-channels are allowed to increase*. Intuitively, at the low rate of primary packets arrival in a given secondary transmission interval, resulting large active sub-channels for the aggregation leads to the increase in the ergodic capacity of an aggregated channel. Moreover, it is worth mentioning that when ξ_{th} increases (towards 1) for a given $\lambda_p \cdot T_d$, n^* in (18) also does. This indicates that large sub-channels can be exploited for the aggregation for a collision-tolerated secondary system with high ξ_{th} . Similarly, as ξ_{th} decreases for the collision-sensitive system, n^* should *exponentially decreases* for a given $\lambda_p \cdot T_d$.

For example, let us consider the homogeneous channel environments where $\gamma_i = \gamma, \forall i$. Using (9) and (18), the maximum ergodic capacity can be derived as follows

$$C_{total} = \left\lfloor \frac{\log(1-\xi_{th})}{-\lambda_p \cdot T_d} \right\rfloor \log_2 \left(1 + \frac{\gamma}{\left\lfloor \frac{\log(1-\xi_{th})}{-\lambda_p \cdot T_d} \right\rfloor} \right), \quad (19)$$

where recall that $\lfloor \cdot \rfloor$ equals n^* in (18). It can be shown from (19) that C_{total} increases monotonically whatever n^* grows, resulting from either 1) low $\lambda_p \cdot T_d$ or 2) high $\xi_{th} (\leq 1)$. It is also worth mentioning that for a given ξ_{th} , C_{total} can remain at a certain minimum desired level, adjusting T_d to the variations of λ_p . That is, when PU packet arrives more frequently (i.e., with large λ_p), it is desired to reduce T_d ,

TABLE I. SIMULATION PARAMETERS

Parameter	Value
The number of total sub-channels	30
PU packet arrival rate, λ_p	[0.02-0.6] /sec
PU packet service time, $1/\mu_p$	1 sec
SU Transmission interval, T_d	10 ms
Total transmit power, P_{max}	1 (Equal power allocation)
Average SNR, $E[g]/\sigma$	20 dB (Rayleigh fading)
Collision probability threshold, ξ_{th}	0.005
Simulation time	$10^5 T_d$

enabling more frequent spectrum sensing. This is in order to accurately detect the PU traffic arrival, leading C_{total} to the desired level.

VI. SIMULATION RESULTS

In this section, we present simulation results to evaluate the formulated collision probabilities and the proposed optimal aggregation algorithm's performance.

Simulations consist in a sequence of events, arrival or leave of primary packets at each sub-channel and secondary transmission using multiple sub-channels. We consider a total of 30 primary channels. Based on channel status information, a BS will randomly choose n sub-channels considering a given collision probability threshold and communicate with the SU. The frequency of collision events are counted and the simulation results are averaged for the obtained results during $10^5 T_d$. The parameters used for simulation are described in Table I.

Firstly, we validate our formulated collision probability that a SU experiences by using multiple sub-channels, described in Section IV. Fig. 1 presents the collision probabilities $P_{c,n}$ formulated in (14) with respect to various λ_p . For given $n \in \{1, 5, 10\}$, λ_p varies from 0.1 to 0.6 per *sec* and $1/\mu_p$ is set to 1 *sec*. Fig. 1 shows that $P_{c,n}$ increases with n and λ_p for a given T_d and that for different values of n , the simulation results closely follow and thus validate the numerical analysis $P_{c,n}$ calculated by (14).

Secondly, the performance of the optimal spectrum aggregation method considering the collision constraint is evaluated by simulation. We compare the performance of the optimal aggregation scheme with three fixed number of sub-channels

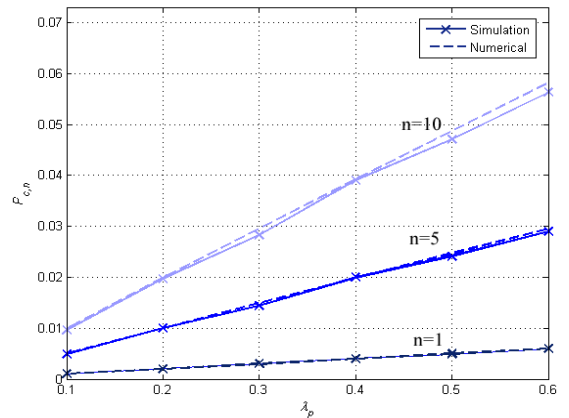
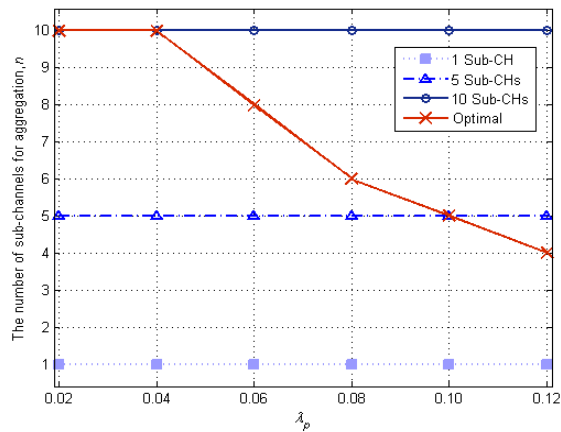
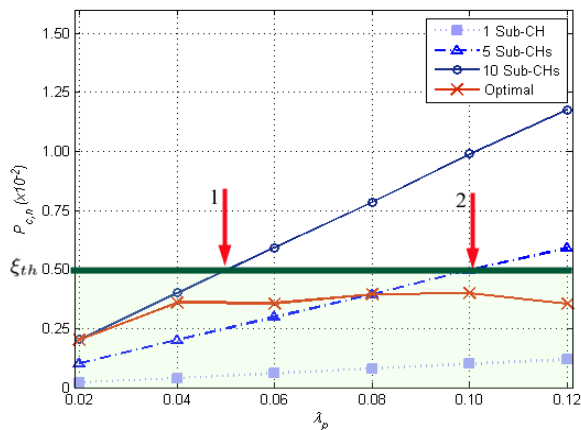


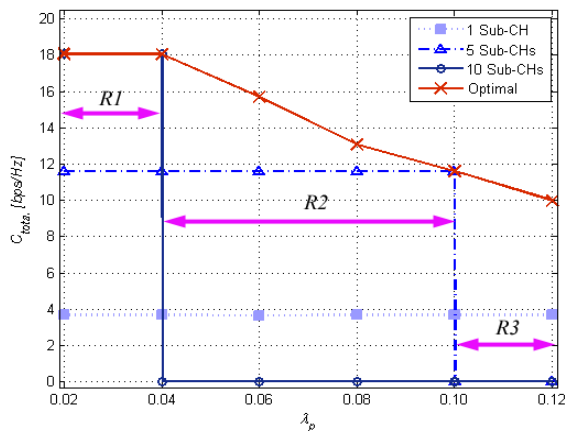
Fig. 1. Comparison of collision probabilities $P_{c,n}$ obtained from simulation and numerical analysis for various λ_p



(a) The number of sub-channels for aggregation



(b) Collision probability of an aggregate channel



(c) Ergodic capacity of an aggregate channel

Fig. 2. Comparison of the proposed aggregation scheme with aggregation of the fixed number of sub-channels

aggregation schemes (i.e., 1, 5, and 10 sub-channels) as shown in Fig. 2. In Fig. 2(a), it is shown that whilst the three fixed aggregation schemes select a given fixed number of sub-channels, the optimal algorithm selects n number of sub-channels adaptively by (18) for given λ_p , T_d and ξ_{th} . In Fig. 2(b), the collision probability $P_{c,n}$ for four schemes is depicted for varying λ_p . Whilst ξ_{th} is set to 0.005 in this scenario, the

valid scheme should generate $P_{c,n}$ lower than ξ_{th} for a given λ_p and T_d . In the case of the “1 Sub-CH” allocation scheme, its $P_{c,n}$ keeps lower value than ξ_{th} for various λ_p . However, for the “10 Sub-CHs” and “5 Sub-CHs” aggregation schemes, their $P_{c,n}$ exceeds ξ_{th} at the points of arrow marked by 1 and 2, respectively. Since we consider the discrete range of λ_p (from 0.02 to 0.12 with an interval of 0.02), under the collision probability constraint, the “10 Sub-CHs” aggregation scheme only remains valid for $\lambda_p = [0.02 - 0.04]$ and the “5 Sub-CHs” aggregation scheme remains valid only for $\lambda_p = [0.02 - 0.10]$. In the case of the optimal aggregation scheme, the $P_{c,n}$ for all λ_p is kept under the ξ_{th} while it changes the number of sub-channels for aggregation. Fig. 2(c) shows the ergodic capacity of the aggregate channel, C_{total} . For region 1 (marked by R1) of $\lambda_p = [0.02 - 0.04]$, the “10 Sub-CHs” aggregation scheme achieves the highest C_{total} among the fixed number of sub-channel aggregation schemes. Due to the $P_{c,n}$ larger than ξ_{th} , for the region of λ_p larger than 0.04, the “10 Sub-CHs” aggregation scheme is no longer valid. The “5 Sub-CHs” aggregation scheme shows better C_{total} than the “1 Sub-CH” allocation scheme for region 2 (marked by R2) of $\lambda_p = (0.04 - 0.10]$. However, for λ_p larger than 0.10, its $P_{c,n}$ becomes larger than ξ_{th} . In region 3 (marked by R3) of $\lambda_p = (0.10 - 0.12]$, only the “1 Sub-CH” allocation scheme can be utilized. While the aggregation algorithms using a fixed number of sub-channel can be utilized within the limited range of λ_p , the optimal algorithm shows the highest C_{total} among four schemes considered for all regions, R1, R2 and R3.

VII. CONCLUSION

We have considered the dynamic spectrum aggregation for the collision-sensitive secondary system so that the ergodic capacity is maximized under the collision probability constraint. It was found that the ergodic capacity by the spectrum aggregation increases with the number of sub-channels activated for aggregation. The collision probability of the aggregate channel was analysed as the function of the PU packet arrival rate, SU transmission interval, and the number of sub-channels activated for the spectrum aggregation. Our analytical and simulation results have clearly shown that the number of sub-channels activated for the spectrum aggregation must exponentially decrease in the highly collision-sensitive use case when decreasing the collision probability threshold. Also, it was found that, for a given collision probability threshold, lower PU packet arrival rate and shorter SU transmission interval can allow a SU to aggregate more sub-channels leading to higher data rates. Based on these findings, the optimal spectrum aggregation strategy has been proposed to achieve the highest ergodic capacity under the given collision probability constraint. Through adaptive selection of number of sub-channels, the proposed optimal scheme was shown to outperform the conventional spectrum aggregation the fixed number of sub-channels. Collisions between primary and secondary transmission can also occur due to noisy estimation of idle channels. Analysis of the impacts of such collisions will be investigated as part of our future work.

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