

Node Credit Based Efficient Flooding (NCBEF) Method for Mobile Ad-hoc Networks

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Abstract

Effective routing is one important research issue for mobile ad-hoc network where gadgets have confined resources. Due to non-foreseeing network topology, frequent link failure occurred that drives route break, partitions and high routing overhead issue. These issues result superfluous resource wastage and diminish network life. Existing reactive routing algorithm helps nodes to discover route via broadcast of route request to all these reachable nodes. Receiving node of route request further forward to all its neighbours' nodes that lead flooding. Here broadcasting of route request results several issues such as wastage of resources, congestions, redundancy of route request and broadcast storm problem. Though to maximize resource utilization and minimize broadcast storm problem, efficient broadcast method in routing algorithm is required. Numerous contributions have done towards efficient broadcasting which concentrate different aspects. In this paper, NCBEF method is presented that motive was to do efficient flooding. NCBEF method enables node to broadcast route request to some neighbours instead of all. To do this, NCBEF method, determine high creditable neighbours during hello signal transmission. Few hypotheses were formulated to evaluate and test NCBEF method. Hypotheses were evaluated by simulation of NCBEF method in network simulator software on the basis of certain criteria.

Keywords: Wireless Networks, Mobile Ad-hoc Networks, Routing, Broadcasting, Flooding, Broadcast Storm Problem, NCBEF.

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1. Introduction

As correspondence bearer point of view, network

classifies in wired and wireless form. Further wireless network likewise had some expertise in

foundation less on the interest of uses and conditions. Figure 1 presents classification of network in in setting of correspondence bearer. In the figure coloured arc highlight the concern which will be discussed in this work.

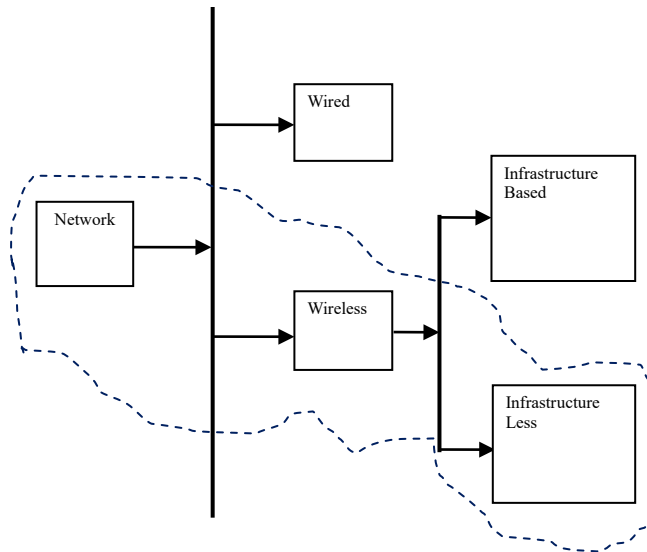


Figure 1. Category of networks

Mobile Ad-hoc Network (MANET) is an example infrastructure less classification where every gadget or node achieves its correspondence paying little respect to pre-set up arrangement. Figure 2 indicates scenario of mobile ad-hoc network that contains four nodes numbered 1 to 4. Here node 1 can convey to node 2 or 3 or 4 and the other way around in light of the fact that they have a place inclusion territory of node1.

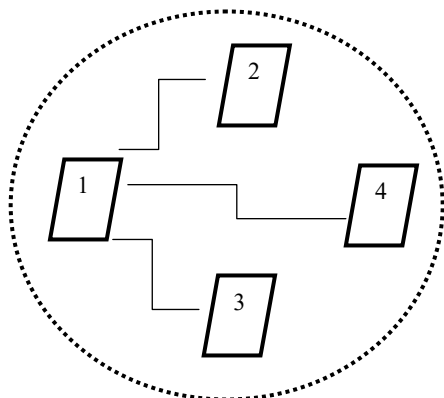


Figure 2. MANET Scenario

In MANET, every node is fit to process routing, route maintenance and re-configuration. Contrarily every node has constrained assets like inclusion territory, battery reinforcement and transmission capacity. Restricted asset limitations property may lead system segments; effective routing and life channel of node also network issues. Customary routing algorithms utilize in assortment of routing conventions which concentrates calculation overhead, accuracy and proficiency. Reactive routing protocol is one of routing protocol types. It works request based methods on the off chance that a node needs send data to other then first it updates routing information through topology change observation and after that discover route. Reactive protocol results less overhead however drives blockage, repetition of routing packets and pointless utilization of assets issues. The reason of issues come about by reactive routing protocols is flooding activity that utilize in route discovery and maintenance. In any case, effective routing is focal central issue of network in nearness of various limitations. The goal of the research is to address effects of flooding activities of reactive routing protocols and attempts to make efficient flooding methods. Further sections present major of flooding, broadcasting, effects of flooding, sorts of flooding methods and recent contributions.

Background

As discussion happened in previous section, flooding method play vital role in route discovery phase of reactive routing protocols. Flooding deals with routing packets to assists routing protocol for route findings. The working of flooding, flooding impacts and types have presented here.

1.1 Flooding

Flooding initiated via broadcasting of routing packets by a node that require route for communication. At this moment receiving node rebroadcast received packets to its covering nodes. Further broadcasting of routing packets spreads throughout the network that results flooding. This method enables routing protocols to tell a node for receiving routing packets and rebroadcast further [1]. Simple scenario of flooding presents in figure 3.

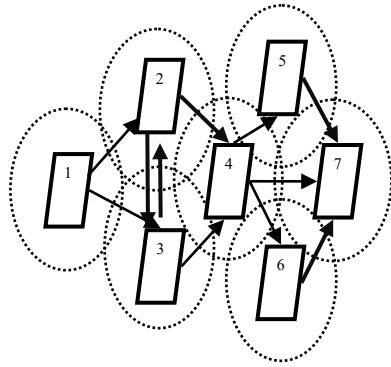


Figure 3. Network Scenarios for Flooding

In above figure, network with seven devices is illustrate where each device labelled with unique number like device 1, 2, 3, 4 etc. Here device begin forwarding of routing messages to its neighbours 2 and 3. Each neighbour receives messages and further forwards to its successor. When device receive same message from all neighbours then process first one and discard others. For example, device 1 receive message from device 1 and 3 then it processes message of device 1 instead of device 3. Likewise, same phenomena follow by all the devices which results flooding.

1.2 Impacts

With the presence of finite resource constraints of devices, ad-hoc network becomes more concentrating field of study and research. Flooding resulted replications of the same packets where each device receives multiple copies of packets which lead numerous problems such as wastage of resource, congestion and many more.

In case of large network, flooding results unnecessary transmission of same routing messages which leads wastage of bandwidth and power, congestions and collisions. Replications of message occurred bandwidth choke and congestion issues. Below figure illustrates discussed issue.

At another site, each device receives messages from those preceding devices and further forwards to those successor devices which are reachable. If network is large then this procedure results storm of messages named broadcast storm problem which impact network functioning or denying services. Figure 4 shows the scenario of storm problem where each device forward to all its successor devices.

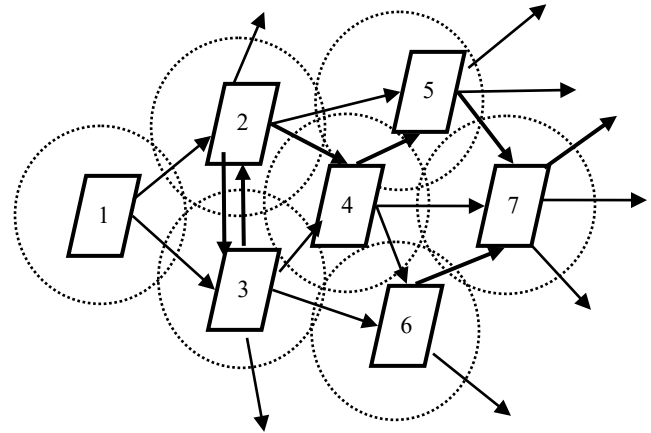


Figure 4. Broadcast Storm Problem

1.3 Flooding Classification

For the reactive routing procedure, flooding is pulse which deal with route discovery in ad-hoc networks. Numerous routing protocols employs flooding like

DSR [2], AODV [3], LAR [4], ZRP [5] etc. Every routing protocol is responsible for route discovery, repair and topology update. Thus, appropriate

flooding methods are required to accomplish routing with less overhead [6, 7]. Various category of flooding method depicted in figure 5.

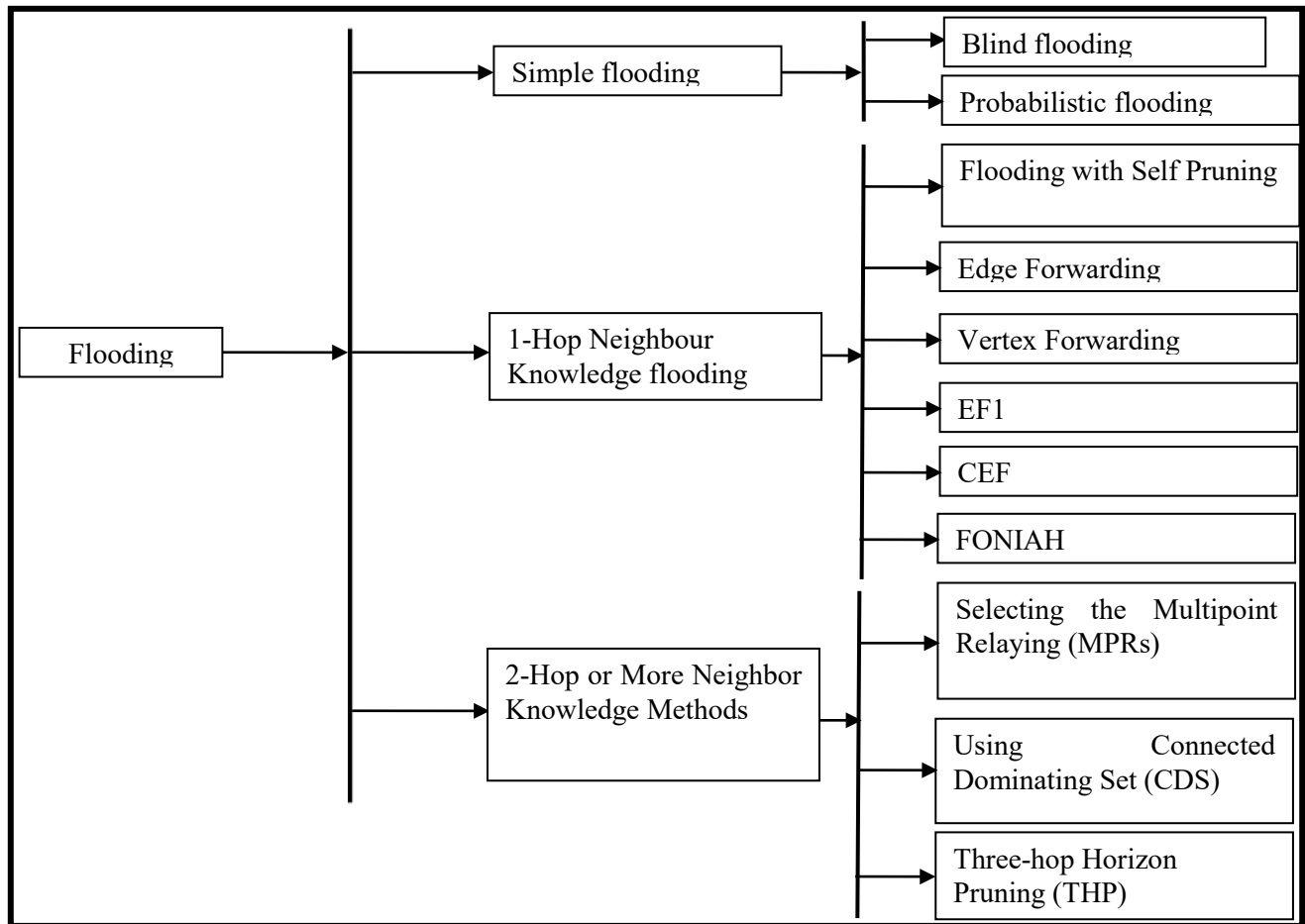


Figure 5. Various category of Flooding

Simple Flooding

This method defines common properties between all the devices of the network where each one forward routing messages without knowing of neighbours. Following sub-category comes in this.

Blind flooding - Blind flooding method alternatively known as pure flooding where every device forwards broadcasting message during route discovery phase [8] [9]. In this, when a device receive forwarding message then it is require to forward further.

Probabilistic flooding - Earlier method resulted

broadcast storm problem which leads wastage of network resources. To sort out this problem, a method suggested that referred as probabilistic [9]. This was employed concept of probabilistic rebroadcast. Therefore, it categorized in counter based, distance based [9], location based [10] and cluster based [11]. This scheme advantages control of redundant transmission of packets over blind method [12].

Hop Neighbor Knowledge Methods

This category method offers knowledge of neighbor's information for every device of the network. Every device gathers and shares its 1-hop neighbor knowledge via exchanging of HELLO message [13].

To choose forwarding nodes, two alternatives are possible. First is, sender node deciding neighbors list for the further forwarding which named sender based. Another one is receiving node decides list of neighbors for further broadcasting which named as receiver based. The schemes expressed in [14] [15] [16] are sender-based, while the schemes proposed in [17] [18] [19] [20] [21] [22] [23] are receiver-based. Various flooding approach [24] [25] [26] falls in 1-hop neighbor knowledge method that ensure about entire message delivery.

2-Hop or More Neighbor Knowledge Methods

In this method, it has assumed that every device contains the record of two-hop neighbors. The one-hop neighbor knowledge gathers by the use of HELLO message in MAC sub-layer of DLL layer. To keep up-to date information about topology, HELLO messages periodically exchanged between devices.

Relevant Work

In preceding section, basic flooding operation and its traditional methods have presented. In few decades, more efforts have putted to contribute resourceful flooding. Here, individual were focused on different parameters and aspects like 0-hop, 1-hop and k-hop neighbour knowledge information. Following section highlight few contributions for the fruitful flooding. Authors in [24] expressed that a one-hop flooding approach accomplishes complete when each device like I the neighbor's scope of I can be come to by $F(I)$, where $F(I)$ is characterized as a subset of I 's neighbors that are chosen to forward the flooding message ($F(I)$ incorporates I itself).The neighbor's

region of a node I is the entire shadow zone appeared in Figure 6, where 2, 3, 4 is the entire neighbors of I 's. Figure 6 shows Neighbor's area of node I .

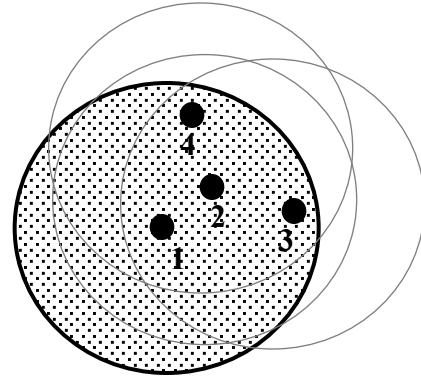


Figure 6. Neighbour's area of node

Another method was advised named as *Flooding with Self-pruning* [27] in which 1-hop neighbour information used. Accordingly, to this method, device forward message by enclosing complete list of 1-hop neighbours. At the receiving site, each node who receive flooding message compare its neighbours with list enclosed in message. Ultimately device either discard forwarding message when it's all 1-hop neighbours is available in the list, or forwards the message further.

Self pruning leads redundant transmission to exchange neighborhood information. To control this, every node maintains received adjacent neighbor records in memory [28].

Further, to offer more efficient flooding an Edge Forwarding [29] method was suggested that uses knowledge of 1-hop neighbor. This method effort to control flooding traffic through geographical knowledge results limited broadcasting within range of each device.

In this method, communication area of node is divided into six uniform subareas. Suppose node I

communication range is divided in six subareas as I_1 , I_2 , I_3 , I_4 , I_5 , I_6 , illustrate in below figure.

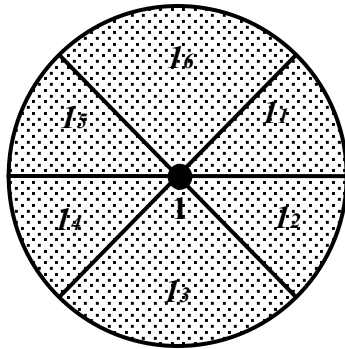


Figure 7. Transmission range of node 1

Authors [19] suggested another scheme named *vertex forwarding* advantaged to routing protocols by reducing traffic with the help neighbor information. Decision made on forwarding node set by the transmission range splitting as the Edge Forwarding [30].

Accordingly, EF1 [25], authors represented network in the form of graph where node was defined vertex that indicated by unique number i. e. v and communication coverage of node was defined as edge that indicated by e .

Therefore, Xian long Jiao *et al* in [31] advised CEF method to optimize performance of routing protocols use of directional antenna. Initially authors expected every node offered inclusive a single-beam directional antenna in EFDA [32], and afterward proposed CEF which is a calculation enhanced EFDA.

Authors [32] stated that for the ad-hoc networks directional antenna was benefited to optimize the performance.

A hybrid flooding strategy was additionally proposed which consolidates "Neighbor learning based flooding and Area-based flooding called

FONIAH (Flooding in light of One-bounce Neighbor Information and Adaptive Holding)" [33].

Further-more an approach advised in which every device aware about the network topography of 2-hop radius. Inappropriately, finding minimum set of devices was NP-hard. Therefore, Amir Qayyum *et al.* [34] suggested a heuristic method to determine the solution named Multipoint Relays (MPR).

Therefore, three-hop horizon pruning (THP) algorithm was the first heuristic to take into account 3-hop information in the selection of relay nodes for broadcast packets [35].

To achieve efficiency, *selective rebroadcasting approach* is further extended in another one referred as MISTRAL [36]. This resulted compensation between overhead and end delivery of messages. Forward error correction technique was adapted to accomplish compensation.

Moreover, another method was designed that known as ASTRAL [37]. A forward and backward packet recovery method was used here. Many norms presented regarding of network topography, coverage of nodes, mobility and MAC protocols.

Further, "*An Efficient Reliable one-hop broadcasting (EROB)*" [38] came in picture which ensures about delivery of message to nodes in the coverage of source node. It leads local one hop broadcasting. It was deal with both data as well as control packets. Data was transmitted via encapsulating in data packets and efficiency enhanced with the help of control packets.

For offering decision capability to nodes for rebroadcasting or discarding receiving packets, a method [39] was advised. In these nodes made

decision on the basis of amount and timestamps of packets.

Global Queue Pruning [40] method was presented that restricts the overhead of the transmission and to enable delivery of packets to each node. It considered virtual topology where few nodes work as forwarders.

At the last but not least, *Selective Epidemic Broadcast Algorithm* [41] method devised to control storm problem of flooding. It was also enhanced rate of data spreading via a new *BSSA based on Selective Epidemic Broadcast Algorithm (SEB)*. The selection of nodes to broadcast packets depends on passive acknowledgement from the neighbours.

2. Proposed Methodology

Apart from discussed all flooding methods, to enable efficient flooding for routing protocols by concentrating limited resource constraints a method is proposed. Proposed method known as *Node Credit Based Efficient Flooding (NCBEF)*. NCBEF method uses one hop neighbour information concept to broadcast routing packets for route discovery. NCBEF assists sender or any other node to select some appropriate neighbours instead of all for further broadcasting. According to NCBEF, sender or any node select high creditable neighbours to broadcast routing packets during routing. The creditability of node is defined by credit (C) which is determined by weight calculation. Therefore, to determine a credit coverage area, stay time and power backup of node have considered. However, Credit (C) is factorized in c1, c2 and c3 with the respect of coverage area, stay time, power backup. The value of c1, c2 and c3 is constant and the sum of all is should be 1. A formula to calculate credit (C) of node defines in below equation.

$$C = c1 * CA + c2 * ST + c3 * PB$$

Where:

CA- CA is coverage area of node

ST- ST is stay time of node at any place

PB- PB is power backup of node

c1, c2 and c3 are credit factors

NCBEF functions in in two steps. First, each node of network determines its credit via calculating weight. Then nodes exchange their credit value with the neighbours to maintain record. Nodes credit exchange achieve through periodic exchange of modified hello signal. In general hello signal leads neighbour discovery in wireless communication. Hello signal is broadcasted in form of message that has number of fields. General format of hello message defined below.

For the NCBEF method, hello message modified by defining new field i.e. C_Val in existing format. The modified hello message shows below.

Src. NID	Hello_Int	Ref_Int	VLT	Seq No	Neighbor ID	C_Val	.	.	Options
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In second step, neighbour table is constructed and updated at each node with the help of modified hello signal. Constructed table keep records of every reachable node and its credit value. Though sender or any broadcasting node use neighbour table to selects higher credit neighbour for further broadcasting.

The whole process of NCBEF method is designed in

Src. NID	Hello_Int	Ref_Int	VLT	Seq No	Neighbor ID	.	.	.	Options
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algorithm that represents in pseudo code form.

2.1 Algorithm

NCBEF method achieves its goal via taking input, processing and resulting output. Thus, node coverage area, stay time and power backup takes as input. Credit calculation defined as process and higher credit value neighbors list as output. The complete process illustrates in below pseudo code.

Algorithm: NCBEF algorithm to select appropriate neighbors for broadcasting

1. Function NCBEF (Node [], N);
Input: CA, ST, PB, i, C_Val, NB
Output: Appropriate neighbors list Neighb[]
2. for i=1 to N do
3. C_Val_i=c1*CA_i+c2* ST_i+c3* PB_i;
4. end for
5. for j=1 to NB do
6. Node[j+1] =Node[j]->Send(hello,C_j);
7. Node[j+1]->update(Node[j], C_j);
8. end for
9. for i=1 to NB-1 do
10. for j=1 to NB do
10. if(C_Val_i>C_Val_{j+1}) then
11. Neighb[i]=i;
12. end for
13. end for
14. for i=1 to N do
- 15 return Neighb[i];
16. end for

Designed algorithm of proposed method takes N number of nodes in Node [], coverage area CA, stay time ST and power backup PB as input. Additionally, to assign credit value of node takes C_Val, NB numbers of neighbors are also takes. It results

appropriate number of neighbors of each node with high credits which listed in Neighb[].

Moreover, to evaluate and test NCBEF method few hypotheses have formulated which may be accepts or rejects. Subsequent subsection represents formulated hypotheses.

2.2 Hypothesis Formulation

NCBEF method has proposed to control packet redundancies, collisions rate and enhance network life by efficient utilization of resource. Thus to evaluate the method some hypotheses formulated which are presents in table 1.

Table 1. Formulated hypotheses

S. No.	Hypotheses	Description
1	h ₁	h ₁ define NCBEF method results less routing overhead
2	h ₂	h ₂ present NCBEF method enhance network life
3	h ₃	h ₃ describe NCBEF method results good routing efficiency
4	h ₄	h ₄ tell NCBEF method results enhance throughput

Formulated hypotheses for NCBEF method will

evaluate and analyze with the existing blind flooding method. Though, experiment will perform with the help of simulation software or tools. Next section defines the whole thing about experiment of proposed method.

3. Experiment of NCBEF Method

In previous section, NCBEF method described hypothetically. Here experimental setup and simulation of proposed method presented. NCBEF method is practically simulated in network simulator software i.e. NS-2. To simulate NCBEF method, various network matrices like packet size, traffic type, number of devices etc. are considers. Table 2 presents network characteristics in form of matrices.

Table 2. Simulation matrices

Matrices Name	Value
Mobile Nodes	20, 40, 60, 80, 100
Topography Area	1200 X 1200
Running Time in Seconds	100
Coverage in meters	150
Signal Propagation	Two Ray Ground
Traffic Rate	CBR, 3pkts/s
Data size in bytes	512
Type of connection	TCP

3.1 Simulation Tool (Network Simulator2)

Network Simulator-2 (NS-2) is function evaluation software for the network which is free available and open source. It is event based that designed and

developed for the communication and network relevant research. NS-2 facilitated numerous types and layer functions of the network.

Network Scenario

To simulate NCBEF method, mobile ad-hoc network scenario is assumes which is comprises of several set of mobile nodes such as 20, 40, 60, 100. Every node is mobile in nature and may act as sender, receiver or forwarder during the transmission. Figure 8 show a network scenario of 100 mobile nodes.

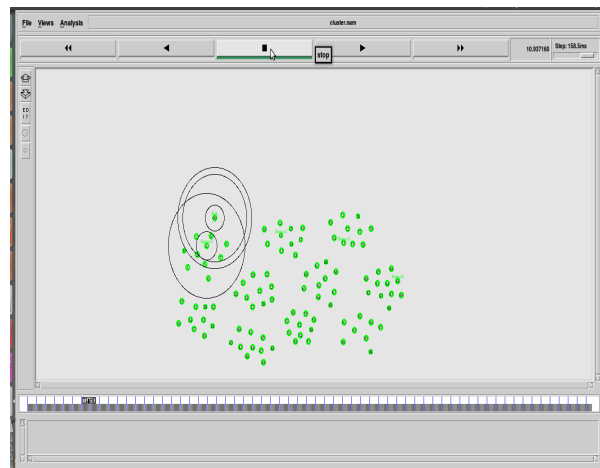


Figure 8. A mobile ad-hoc networks scenario

4. Evaluation of NCBEF Method

NCBEF method is evaluated through test of hypotheses which have been formulated in proposed methodology section. Formulated hypotheses tested on the basis of routing efficiency, routing overhead, throughput and network life. Moreover, NCBEF method is evaluated with comparisons of blind flooding method.

Routing Overhead - Number of routing request employs to route findings and repairs are called routing overhead. The unit is used to measure routing overhead is total number of routing request used in

whole transmission. Figure 9 shows comparative routing overhead of both methods.

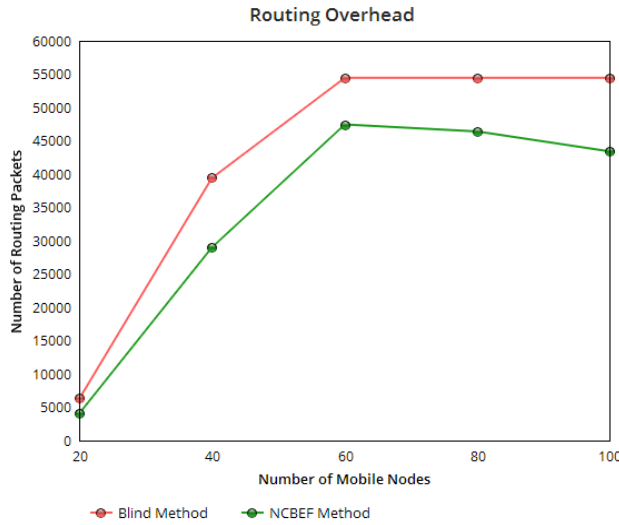
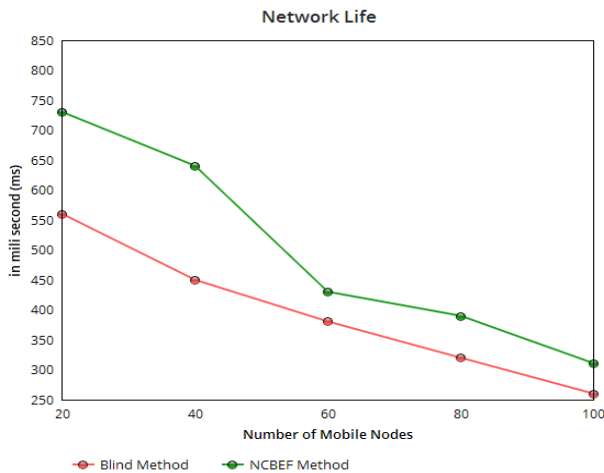


Figure 9. Comparative routing overhead

Network Life - NCBEF method results more network life in time period unit's comparatively blind method. Figure 10 show network life comparison of



both methods.

Figure 10. Comparative network life

Routing Efficiency - Routing Efficiency is characterized as proportion of information transmitted to the aggregate routing requests and information packets. Comparative routing efficiency

of both methods presents in figure 11.

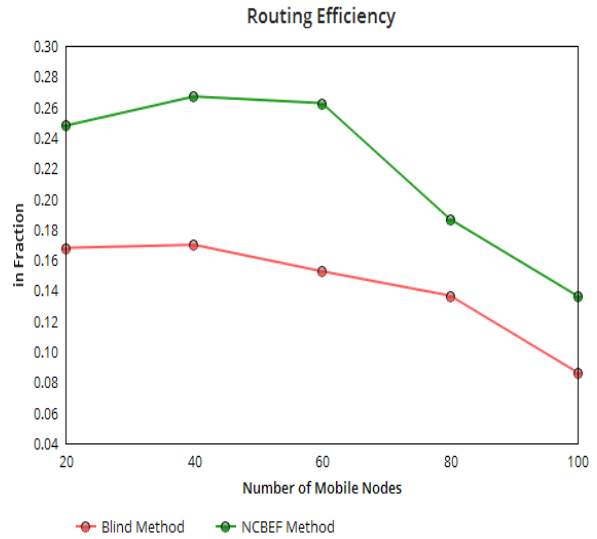


Figure 11. Comparative routing efficiency

Throughput - NCBEF method results more throughput comparatively blind method. Throughput is data units received successfully per unit time at the receiver. Figure 12 show routing efficiency comparison of both methods.

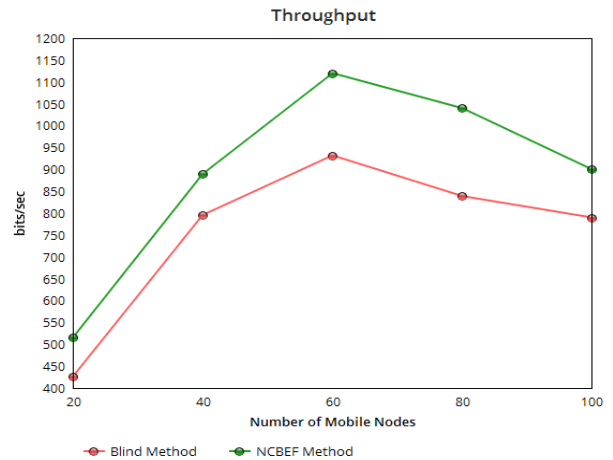


Figure 12. Comparative throughput

4.1 Result Discussions

Accordingly to test formulated hypotheses for evaluation of NCBEF method, more routing

efficiency, network life and less overhead observed. Means, formulated hypotheses are accepted as per results. May be some hypotheses rejects for the same.

5. Conclusion

Due to non-anticipating network topology, frequent link failure occurred that leads route break, partitions and high routing overhead issue. These issues results unnecessary assets wastage and diminish network life. Existing reactive routing algorithm encourages nodes to discover route through communicate of course solicitation to all these reachable nodes. Receiving node of route request further forward to all its neighbours' nodes that lead flooding. Here communicating of course solicitation results a few issues, for example, wastage of assets, blockages, excess of course solicitation and communicate storm issue. However, NCBEF technique empowers hub to communicate course solicitation to certain neighbors rather than all. To do this, NCBEF strategy, decide high respectable neighbors during hello signal transmission. Few hypotheses were formulated to evaluate and test NCBEF method which was accepted.

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